1 Better Distortion with Distribution

There is a well known lowerbound for the distortion of embedding a metric space (X, d) into ℓ_1 .

Theorem 1.1 For any metric space (X, d) on n points, one has

$$(X,d) \stackrel{\Omega(\log n)}{\longleftrightarrow} \ell_1.$$

For ℓ_2 the lowerbound is still $\Omega(\log n)^{-1}$.

Recall that we want to find a $(O(k), (1+\varepsilon)c)$ -outlier embedding into ℓ_2 for any metric space (X,d) which admits a (k,c)-outlier embedding into ℓ_2 . If we can do this deterministically, we actually find an embedding of the outlier points into ℓ_2 with distortion O(k), which contradicts the lowerbound. However, maybe we can do O(k) via embedding into some distribution of ℓ_2 metrics.

Let (X, d) be a finite metric space and let $\mathcal{Y} = \{(Y_1, d_1), \dots (Y_h, d_h)\}$ be a set of metric spaces. Let π be a distribution of embeddings into \mathcal{Y} . The original metric space (X, d) embeds into π with distortion D if there is an r > 0 such that for all $x, y \in X$,

$$r \le \frac{\mathrm{E}_{i \leftarrow \pi}[d_i(\alpha_i(x), \alpha_i(y))]}{d(x, y)} \le Dr.$$

SODA23 paper also embeds (X,d) into distribution. We call this kind of embeddings stochastic embedding.

Example: Random Trees Consider the problem of embedding some finite metric into a tree metric. We can get an O(n) lowerbound via the unit edge length cycle C_n . However, if embedding into distortions is allowed, we can do $O(\log n)$.

Theorem 1.2 (Bartal) Let (X,d) be a metric space on n points with diameter Δ , let DT be the set of tree metrics that dominate d, there is a distribution π on DT such that (X,d) embeds into pi with distortion $O(\log n)$.

2 Stochastic Embedding into ℓ_2

We first ignore the outlier condition and see if stochastic embeddings break the $\Omega(\log n)$ lower-bound.

Theorem 2.1 (Bourgain) For any metric space (X,d) and for any p, there is an embedding of (X,d) into $\ell_p^{O(\log^2 n)}$ with distortion $O(\log n)$.

We want to beat Bourgain's in terms of the expected distortion.

¹https://web.stanford.edu/class/cs369m/cs369mlecture1.pdf